

n -fold Obstinate Filters in Pseudo-Hoop Algebras

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Abstract

In this paper, we introduce the concepts of n -fold obstinate pseudo-hoop and n -fold obstinate filter in pseudo-hoops. Then we investigated these notions and proved some properties of them. Also, we discussed the relationship between n -fold obstinate pseudo-hoop and n -fold obstinate filter and other types of n -fold pseudo-hoops and n -fold filters such as n -fold(positive) implicative filter and n -fold fantastic filter in pseudo-hoops. For example, we proved that any n -fold obstinate filter is a maximal filter. Finally, we obtain a characterization of n -fold obstinate filters in terms of congruences and we show that any n -fold obstinate pseudo-hoop is an n -fold fantastic, n -fold positive implicative, n -fold implicative pseudo-hoop and simple pseudo-hoop.

Keywords : Pseudo-hoop algebra; Filter; n -fold obstinate pseudo-hoop; n -fold obstinate filter.

1 Introduction

Naturally ordered commutative residuated integral monoids (hoop) introduced by B. Bosbach in [5, 6], then studied by J. R. Büchi et al. in [7], a paper never published. Also G. Georgescu, L. Leustean et al. study the pseudo-hoops in [8]. It is well-known that in various logical systems, filters play a fundamental role, filters correspond to sets of provable formulas closed with respect to Modus Ponens. In [10, 12, 14, 16, 17] the authors investigated the notation folding theory to residuated lattices, n -folding fantastic filters and obstinate filters in BL-algebras, general-

ization of integral filters and n -fold integral BL-algebras and n -fold filters of MTL-algebras. In [2], R. A. Borzooei et al., survey the notion of n -fold(implicative, positive implicative and fantastic filters) of pseudo-hoops. They show that if F is an n -fold(implicative, positive implicative and fantastic)filter, then A/F is an n -fold (implicative, positive implicative and fantastic)pseudo-hoops. Also in [15], A. Namdar et al., proposed the obstinate filter in hoops.

In this disquisition, we define and study the notion of n -fold obstinate pseudo-hoop and n -fold obstinate filters in pseudo-hoops and generalization of the corresponding notion in the crisp case. Several properties of n -fold obstinate pseudo-hoop and n -fold obstinate filters are given. We show that F is an n -fold obstinate filter of A if and only if A/F is an n -fold obstinate pseudo-hoop. On the other hands if F is an n -fold obstinate filter of A , then A/F is a local and simple pseudo-hoop. Also, we show that F is an n -fold

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obstinate filter if and only if F is a maximal and n -fold positive implicative filter.

2 Preliminaries

In this section, we recollect some definitions and results which will be used in this paper.

Definition 2.1 [8] *A pseudo-hoop algebra or pseudo-hoop is an algebra $(A, \odot, \rightarrow, \rightsquigarrow, 1)$ of type $(2, 2, 2, 0)$ such that, for all $x, y, z \in A$:*

- (PH1) $x \odot 1 = 1 \odot x = x$,
- (PH2) $x \rightarrow x = x \rightsquigarrow x = 1$,
- (PH3) $(x \odot y) \rightarrow z = x \rightarrow (y \rightarrow z)$,
- (PH4) $(x \odot y) \rightsquigarrow z = y \rightsquigarrow (x \rightsquigarrow z)$,
- (PH5) $(x \rightarrow y) \odot x = (y \rightarrow x) \odot y = x \odot (x \rightsquigarrow y) = y \odot (y \rightsquigarrow x)$.

On pseudo-hoop A , we define $x \leq y$ if and only if $x \rightarrow y = x \rightsquigarrow y = 1$. It is easy to see that \leq is a partial order relation on A . If \odot is commutative (or equivalently $\rightarrow = \rightsquigarrow$), then A is said to be a hoop. A pseudo-hoop A is *bounded* if there is an element $0 \in A$ such that $0 \leq x$, for all $x \in A$. For any $x \in A$, we consider $x^- = x \rightarrow 0$ and $x^\sim = x \rightsquigarrow 0$. An element $x \in A$ is called *atom* if it is a minimal among elements in bounded hoop $A \setminus \{0\}$. Also, element $x \in A$ is called *idempotent* if $x^2 = x$. The *order* of $1 \neq x \in A$, in symbols $ord(x)$ is the smallest $n \in \mathbb{N}$ such that $x^n = 0$. If no such n exists, then $ord(x) = \infty$. (See [8])

Definition 2.2 [8] *For pseudo-hoop A and for any $x, y \in A$, we define $x \vee y = ((x \rightarrow y) \rightsquigarrow y) \wedge ((y \rightarrow x) \rightsquigarrow x) = ((x \rightsquigarrow y) \rightarrow y) \wedge ((y \rightsquigarrow x) \rightarrow x)$. If \vee is the join operation on A , then A is called a pseudo \vee -hoop.*

Proposition 2.1 [8] *In any pseudo-hoop A , the following properties hold, for all $x, y, z \in A$:*

- (i) (A, \leq) is a meet-semilattice with $x \wedge y = (x \rightarrow y) \odot x = x \odot (x \rightsquigarrow y)$,
- (ii) $1 \rightarrow x = x, 1 \rightsquigarrow x = x, x \rightsquigarrow x = 1$,
- (iii) $y \leq x \rightarrow y$ and $y \leq x \rightsquigarrow y$,
- (iv) if $x \leq y$, then $y \rightsquigarrow z \leq x \rightsquigarrow z$ and $y \rightarrow z \leq x \rightarrow z$,
- (v) $x \odot y \leq x, y$ and $x^n \leq x$, for any $n \in \mathbb{N}$,
- (vi) if \vee exists, then $(x \vee y) \rightsquigarrow z = (x \rightsquigarrow z) \wedge (y \rightsquigarrow z)$, $(x \vee y) \rightarrow z = (x \rightarrow z) \wedge (y \rightarrow z)$.

Proposition 2.2 [8] *Let A be a bounded pseudo-hoop. Then the following properties hold, for all $x, y, z \in A$:*

- (i) if $x \leq y$, then $y^\sim \leq x^\sim$ and $y^- \leq x^-$,
- (ii) $(x^n)^- \leq (x^{n+1})^-$ and $(x^n)^\sim \leq (x^{n+1})^\sim$,
- (iii) $0^- = 0^\sim = 1$ and $1^- = 1^\sim = 0$,
- (iv) $x \leq (x^-)^\sim$ and $x \leq (x^\sim)^-$,
- (v) $x \odot x^- = x \odot x^\sim = 0$,
- (vi) $x^- \leq x \rightarrow y$ and $x^\sim \leq x \rightsquigarrow y$.

Definition 2.3 [8] *Let A be a pseudo-hoop. A non-empty subset F of A is called a filter of A if,*

- (F1) $x \in F$ and $x \leq y$, then $y \in F$, for any $x, y \in A$,
- (F2) $x \odot y \in F$, for any $x, y \in F$.

Clearly, $1 \in F$, for all filters of A . A filter F of A is called a *proper filter* if $F \neq A$. It is easy to see that, if A is a bounded pseudo-hoop, then a filter is proper if and only if it is not containing 0. The set of all filters of A denoted by $\mathcal{F}(A)$.

Proposition 2.3 [8] *Let A be a pseudo-hoop. If F is a non-empty subset of pseudo-hoop A such that $1 \in F$, then the following statements are equivalent, for any $x, y \in A$:*

- (i) F is a filter,
- (ii) if $x, x \rightarrow y \in F$, then $y \in F$,
- (iii) if $x, x \rightsquigarrow y \in F$, then $y \in F$.

Notation: It is easy to see that the intersection of all filters of pseudo-hoop A is a filter. Hence, for any $B \subseteq A, \bigcap_{B \subseteq F \in \mathcal{F}(A)} F$ is a filter and denoted by $[B]$ and we called *generated filter by B* .

Theorem 2.1 [8] *Let $x \in A$. Then $[x] = \{a \in A \mid x^n \leq a, \text{ for some } n \geq 1\}, F(x) = [F \cup \{x\}] = \{t \mid t \geq f \odot x^n \text{ for } f \in F, n \in \mathbb{N}\}$ and $[F \cup G] = \{a \in A \mid a \geq f \odot g \text{ for } f \in F, g \in G\}$, for any $F, G \in \mathcal{F}(A)$.*

Definition 2.4 [8] *A filter F of pseudo-hoop A is called a normal filter if $x \rightarrow y \in F$ if and only if $x \rightsquigarrow y \in F$, for all $x, y \in A$.*

Definition 2.5 [8] *A proper filter F of a pseudo \vee -hoop A is called a prime filter of A if $x \vee y \in F$, then $x \in F$ or $y \in F$, for any $x, y \in A$.*

A maximal filter of pseudo-hoop A is a proper filter M of A that is not included in any other proper filters of A . $Max(A)$ is the set of all maximal filters of A .

Proposition 2.4 [8] *Let A be a pseudo-hoop and F be a non-empty subset of pseudo-hoop A . Then the following conditions are equivalent, for any $x \in A$:*

- (i) F is a maximal filter,
- (ii) $x \notin F$ if and only if $(x^n)^-, (x^n)^\sim \in F$, for some $n \in \mathbb{N}$.

Proposition 2.5 [3] *Let A be a bounded \vee -hoop. Then every maximal filter of A is a prime filter.*

Definition 2.6 [2] *Let F be a subset of A such that $1 \in F$. Then for any $x, y, z \in A$:*

- (i) F is called an n -fold positive implicative filter of A , if $x^n \rightarrow (y \rightarrow z) \in F$ and $x^n \rightsquigarrow y \in F$, then $x^n \rightarrow z \in F$. Also, if $x^n \rightsquigarrow (y \rightsquigarrow z) \in F$ and $x^n \rightarrow y \in F$, then $x^n \rightsquigarrow z \in F$.
- (ii) F is called an n -fold implicative filter of A , if $x \rightarrow ((y^n \rightarrow z) \rightsquigarrow y) \in F$ and $x \in F$, then $y \in F$. Also, if $x \rightsquigarrow ((y^n \rightsquigarrow z) \rightarrow y) \in F$ and $x \in F$, then $y \in F$.
- (iii) F is called an n -fold fantastic filter of A , if $z \rightarrow (y \rightarrow x) \in F$ and $z \in F$, then $((x^n \rightarrow y) \rightsquigarrow y) \rightarrow x \in F$. Also, if $z \rightsquigarrow (y \rightsquigarrow x) \in F$ and $z \in F$, then $((x^n \rightsquigarrow y) \rightarrow y) \rightsquigarrow x \in F$.

Definition 2.7 [8] *Let A and B be two bounded pseudo-hoops. A map $f : A \rightarrow B$ is called a pseudo-hoop homomorphism if and only if for all $x, y \in A$, $f(0) = 0$, $f(1) = 1$, $f(x \odot y) = f(x) \odot f(y)$, $f(x \rightarrow y) = f(x) \rightarrow f(y)$ and $f(x \rightsquigarrow y) = f(x) \rightsquigarrow f(y)$.*

The set of all pseudo-hoop homomorphism from A to B is shown by $Hom(A, B)$.

Definition 2.8 [8] *Let A be a pseudo-hoop. Then A is called:*

- (i) n -fold positive implicative pseudo-hoop, if $x^{n+1} = x^n$, for all $x \in A$.
- (ii) n -fold implicative pseudo-hoop, if $(x^n \rightarrow 0) \rightsquigarrow x = x$ and $(x^n \rightsquigarrow 0) \rightarrow x = x$, for all

$x \in A$.

- (iii) n -fold fantastic pseudo-hoop, if $((x^n \rightarrow y) \rightsquigarrow y) \rightarrow x = y \rightarrow x$ and $((x^n \rightsquigarrow y) \rightarrow y) \rightsquigarrow x = y \rightsquigarrow x$, for all $x, y \in A$.
- (iv) local pseudo-hoop, if $ord(x) < \infty$ or $ord(x^-) < \infty$ or $ord(x^\sim) < \infty$, for all $x \in A$.
- (v) simple pseudo-hoop, if A is non-trivial and $\{1\}$ is its only proper filter.
- (vi) cancellative pseudo-hoop, if the monoid $(A, \odot, 1)$ is cancellative if and only if $b \rightarrow (a \odot b) = a$ and $b \rightsquigarrow (a \odot b) = a$ if and only if $c \odot a = c \odot b$, then $a = b$, for any $a, b, c \in A$.

Notation: From now on, we let $(A, \odot, \rightarrow, \rightsquigarrow, 0, 1)$ or A be a bounded pseudo-hoop, unless otherwise state.

3 n-fold obstinate pseudo-hoops and n-fold obstinate filters in pseudo-hoops

In this section, we introduce the notion of n -fold obstinate pseudo-hoop and n -fold obstinate filter in pseudo-hoop and investigate some properties of them.

Definition 3.1 *A is called an n -fold obstinate pseudo-hoop if, for all $x \neq 1$, $x^n = 0$.*

Example 3.1 (i) *Let $(A = \{0, a, b, 1\}, \leq)$ be a chain that is $0 < a < b < 1$. Define the operations \odot, \rightarrow and \rightsquigarrow on A as follows:*

$\rightarrow, \rightsquigarrow$	0	a	b	1
0	1	1	1	1
a	b	1	1	1
b	a	b	1	1
1	0	a	b	1

\odot	0	a	b	1
0	0	0	0	0
a	0	0	0	a
b	0	0	a	b
1	0	a	b	1

Then $(A, \odot, \rightarrow, \rightsquigarrow, 1, 0)$ is a bounded pseudo-hoop and A is an 3-fold obstinate pseudo-hoop.

But it is not an 2-fold obstinate pseudo-hoop, because $b^2 \neq 0$.

(ii) [9]. Let $NS[0, 1]$, (non-standard interval $[0, 1]$) be the ordered set whose elements are pairs (a, b) such that $a = 0$ and $0 \leq b$ or $0 < a < 1$ and b arbitrary or $a = 1$ and $b \leq 0$ (b running on real set). The ordering is lexicographic: $(a, b) \leq (c, d)$ if and only if $a < c$ or $(a = c$ and $b \leq d)$. The ordered set $NS[0, 1]$ endowed with the operations: $(a, b) \odot (c, d) =$

$$\max \left((0, 0), \left(\frac{1}{2}(a + c - 1 + ac), \frac{b(c + 1)}{2} \right) \right)$$

If $(a, b) \leq (c, d)$, then $(a, b) \rightarrow (c, d) = 1$, otherwise $(a, b) \rightarrow (c, d) = \left(\frac{2c-a+1}{1+a}, \frac{2d-2b}{1+a} \right)$.

Also, if $(a, b) \leq (c, d)$, then $(a, b) \rightsquigarrow (c, d) = 1$, otherwise $(a, b) \rightsquigarrow (c, d) = \left(\frac{2c-a+1}{1+a}, \frac{-b(c+1)}{1+a} + d \right)$.

Then $(NS[0, 1], \odot, \rightarrow, \rightsquigarrow, 0, 1)$ is a bounded pseudo-hoop. But it is not an n -fold obstinate pseudo-hoop, because $(1, b) \odot (1, b) = \max((0, 0), (1, b)) = (1, b)$, for $b \leq 0$.

(iii) [1] Let $A = [0, \frac{1}{2}] \cup \{1\}$ and operations \odot, \rightarrow and \rightsquigarrow are defined by, $x \odot y = \max(0, x + y - 1)$, and if $x \leq y$, then $x \rightarrow y = 1$, otherwise $x \rightarrow y = \min(1 - x + y, 1)$.

Then $(A, \odot, \rightarrow, 1, 0)$ is an 2-fold obstinate pseudo-hoop.

Proposition 3.1 *If A is an n -fold obstinate pseudo-hoop, then A is an $(n + 1)$ -fold obstinate pseudo-hoop.*

Proof. Let A be an n -fold obstinate pseudo-hoop. Then $x^n = 0$, for any $x \in A \setminus \{1\}$. By Proposition 2.1(v), $x^{n+1} \leq x^n$. Hence, $x^{n+1} = 0$, for any $x \in A \setminus \{1\}$ and so A is an $(n + 1)$ -fold obstinate pseudo-hoop.

Corollary 3.1 *Any n -fold obstinate pseudo-hoop is an $(n + k)$ -fold obstinate pseudo-hoop, for all $k \geq 1$.*

Proposition 3.2 *If A is an n -fold obstinate pseudo-hoop, then A is not a cancellative pseudo-hoop.*

Proof. Let A be a cancellative pseudo-hoop, by the contrary. Then $x^{n+1} = x^n = 0$. Hence $x^n \odot$

$x = x^n \odot 1 = 0$. Thus $x = 1 = 0$, which is a contradiction. Therefore, A is not a cancellative pseudo-hoop.

Proposition 3.3 *If A does not have idempotent element except $\{0, 1\}$ and $\mathcal{A}(M)$ is the set of all atoms of A , then $\mathcal{A}(M) \cup \{1\}$ is an n -fold obstinate pseudo-hoop.*

Proof. If $x \in \mathcal{A}(M)$, then x is an atom and is not idempotent element of A . Thus $x^2 \neq x$. By Proposition 2.1(v), $x^n = x^2 = 0$.

Proposition 3.4 *If A is an n -fold obstinate pseudo-hoop, then A does not have idempotent element except $0, 1$.*

Proof. Let $0 \neq x$ be an idempotent element of A . Then $x^2 = x$. Since A is an n -fold obstinate pseudo-hoop, $0 = x^n = x$, which is a contradiction.

Notation: For any $x \in A$, we consider $m_x = \text{ord}(x) - 1$, so $x^{m_x} \neq 0$.

Proposition 3.5 *Let A be an n -fold obstinate pseudo-hoop. Then x^{m_x} is an atom for any $0, 1 \neq x \in A$ and $m_x \in \mathbb{N}$.*

Proof. Let $x \in A$. Then $x^n = 0$ and $0 = x^n \leq x^{n-1} \leq x^{n-2} \leq \dots \leq x$. If $t = \text{ord}(x)$, then $x^{t-1} \neq 0$. So for $m_x = t - 1$, x^{m_x} is an atom.

Definition 3.2 *A proper filter F of A is called an n -fold obstinate filter if for all $x, y \notin F$, then $x^n \rightarrow y, y^n \rightarrow x \in F$ and $x^n \rightsquigarrow y, y^n \rightsquigarrow x \in F$, for $n \in \mathbb{N}$.*

Example 3.2 *In Example 3.1(i), $F = \{1\}$ is an 3-fold obstinate filter but since $b^2 \rightarrow 0 = a \rightarrow 0 = b \notin F$, F is not an 2-fold obstinate filter of A .*

Proposition 3.6 *Let F be a proper filter of A . Then the following statements are equivalent:*

- (i) F is an n -fold obstinate filter of A ,
- (ii) $x \in F$ or $(x^n)^-, (x^n)^\sim \in F$, for all $x \in A$.

Proof. (i) \Rightarrow (ii) Suppose F is an n -fold obstinate filter and $x \notin F$. Since F is a proper filter and A is bounded, $0 \notin F$. Then $(x^n)^- = x^n \rightarrow 0 \in F$ and $(x^n)^\sim = x^n \rightsquigarrow 0 \in F$.

(ii) \Rightarrow (i) Let $x, y \notin F$. Then by assumption, $(x^n)^-, (x^n)^\sim, (y^n)^-, (y^n)^\sim \in F$. Thus,

by Proposition 2.2(vi), $(x^n)^- \leq x^n \rightarrow y$ and $(y^n)^- \leq y^n \rightarrow x$. Since F is a filter, by (F1), $x^n \rightarrow y \in F$ and $y^n \rightarrow x \in F$. The proof of other case is similar. Therefore, F is an n -fold obstinate filter of A .

Corollary 3.2 F is an n -fold obstinate filter of A if and only if $x \notin F$ implies $((x^n)^-)^m, ((x^n)^{\sim})^m \in F$, for all $m \in \mathbb{N}$ and $x \in A$.

Proposition 3.7 Let F be an n -fold obstinate filter of A . Then the following conditions hold:

- (i) for all $0, 1 \neq x \in A$, $x^n \rightarrow (x^n)^-$, $x^n \rightarrow (x^n)^{\sim} \in F$ or $(x^n)^- \rightarrow x^n$, $(x^n)^{\sim} \rightarrow x^n \in F$,
- (ii) for all $x \in A$, $((x^n)^-)^{\sim} \rightarrow x^n$, $((x^n)^{\sim})^- \rightsquigarrow x^n \in F$,
- (iii) for all $x \notin F$, and for any $y \leq x^n$, then $y^-, y^{\sim} \in F$,
- (iv) for all $x \in A$, $x^n \rightarrow x^{2n}$, $x^n \rightsquigarrow x^{2n} \in F$.

Proof. (i) Let $x \in F$. Then by Proposition 2.1(iii), $x^n \leq (x^n)^- \rightarrow x^n$. Since F is a filter, by (F1), $(x^n)^- \rightarrow x^n \in F$. If $x \notin F$, then by Proposition 3.6(ii), $(x^n)^- \in F$. By Proposition 2.1(iii), $(x^n)^- \leq x^n \rightarrow (x^n)^-$, and so $x^n \rightarrow (x^n)^- \in F$. The proof of other cases is similar.

(ii) We consider the following cases:

Case 1: If $x \in F$, then by Proposition 2.1(iii), $x^n \leq ((x^n)^-)^{\sim} \rightarrow x^n$. Since F is a filter, by (F1), $((x^n)^-)^{\sim} \rightarrow x^n \in F$.

Case 2: If $x \notin F$, then by Proposition 3.6(ii), $(x^n)^- \in F$. By Proposition 2.2(iv) and (vi), $(x^n)^- \leq (((x^n)^-)^{\sim})^- \leq ((x^n)^-)^{\sim} \rightarrow x$, and so by (F1), $((x^n)^-)^{\sim} \rightarrow x \in F$. The proof of other cases is similar, too.

(iii) Let $x \notin F$ and $y \leq x^n$. Then by Proposition 2.2(i), $(x^n)^- \leq y^-$. Since F is an n -fold obstinate filter, by Proposition 3.6(ii), $(x^n)^- \in F$ and by (F1), $y^- \in F$.

(iv) We consider the following cases:

Case 1: If $x \in F$, then by Proposition 2.1(iii), $x^{2n} \leq x^n \rightarrow x^{2n}$. Since F is a filter, by (F1), $x^n \rightarrow x^{2n} \in F$.

Case 2: If $x \notin F$, then by Proposition 3.6(ii), $(x^n)^- \in F$. By Proposition 2.2(iv) and (vi), $(x^n)^- \leq (((x^n)^-)^{\sim})^- \leq ((x^n)^-)^{\sim} \rightarrow x^{2n}$. Also, by Proposition 2.2(iv) and Proposition 2.1(iv),

$x^n \leq ((x^n)^-)^{\sim}$ and $((x^n)^-)^{\sim} \rightarrow x^{2n} \leq x^n \rightarrow x^{2n}$. Therefore, by (F1), $x^n \rightarrow x^{2n} \in F$.

Proposition 3.8 If F is an n -fold obstinate filter of A , then F is an $(n+k)$ -fold obstinate filter of A , for any $k \in \mathbb{N}$.

Proof. Let $x \notin F$. Then by Proposition 3.6(ii), $(x^n)^- \in F$. By Proposition 2.2(ii), $(x^n)^- \leq (x^{n+1})^-$ and so by (F1), $(x^{n+1})^- \in F$. The proof of other case is similar.

Let $F \in \mathcal{F}(A)$. Define $x \equiv_F y$ if and only if $x \rightarrow y \in F$, $y \rightarrow x \in F$, and $x \rightsquigarrow y \in F$, $y \rightsquigarrow x \in F$ for any $x, y \in A$. Then we can see that \equiv_F is a congruence relation on A . The set of all congruence classes is denoted by A/F , it means $A/F = \{[x] \mid x \in A\}$, where $[x] = \{y \in A \mid x \equiv_F y\}$. Define the operations \odot, \rightarrow and \rightsquigarrow on A/F by $[x] \odot [y] = [x \odot y]$, $[x] \rightarrow [y] = [x \rightarrow y]$ and $[x] \rightsquigarrow [y] = [x \rightsquigarrow y]$. Therefore, $(A/F, \odot, \rightarrow, \rightsquigarrow, [1], [0])$ is a bounded pseudo-hoop with respect to F and $[x] \leq [y]$ if and only if $x \rightarrow y, x \rightsquigarrow y \in F$. (See [8])

Notation: It is easy to show that every obstinate filter of A is an n -fold obstinate filter of A and every 1-fold obstinate filter of A is an obstinate filter of A .

Theorem 3.1 Let F be an 1-fold obstinate filter of A . Then A/F is a Boolean algebra.

Proof. Let $x \in A$. Since F is an 1-fold obstinate filter, by Proposition 3.6(ii), $x \in F$ or $x^-, x^{\sim} \in F$. Then, $[x] = [1]$ or $[x^-] = [x^{\sim}] = [1]$. Hence, $[x] = [1]$ or $[(x^-)^{\sim}] = [0]$. If $[(x^-)^{\sim}] = [0]$, since $[x] \leq [(x^-)^{\sim}]$, then $[x] = [0]$. Therefore, A/F is a Boolean algebra.

Theorem 3.2 F is an n -fold obstinate filter of A if and only if A/F is an n -fold obstinate pseudo-hoop.

Proof. (\Rightarrow) Let F be an n -fold obstinate filter and $x \notin F$. Then $x/F \neq 1/F$. By Proposition 3.6(ii), $(x^n)^- \in F$, thus $(x^n)^-/F = 1/F$. By Proposition 2.2(ii) and (iii), $x^n/F = 0/F$.

(\Leftarrow) Let A/F be an n -fold obstinate pseudo-hoop and $x \notin F$. Then $x^n/F = 0/F$ and by Proposition 2.2(iii), $(x^n)^-/F = 1/F$. Hence $(x^n)^- \in F$.

By Proposition 3.6(ii), F is an n -fold obstinate filter of A .

Proposition 3.9 *Let F and G be two filters of A such that $F \subseteq G$. If F is an n -fold obstinate filter of A , then G is an n -fold obstinate filter, too.*

Proof. Let F and G be two filters of A such that $F \subseteq G$ and F be an n -fold obstinate filter of A . Suppose $x \notin G$. Then $x \notin F$. Since F is an n -fold obstinate filter, by Proposition 3.6(ii), $(x^n)^-, (x^n)^\sim \in F$. Hence $(x^n)^-, (x^n)^\sim \in G$ and G is an n -fold obstinate filter of A .

Proposition 3.10 *Let F be an n -fold obstinate filter of A . Then:*

- (i) $(x \odot y)^- \in F$, implies $(x^n)^- \in F$ or $(y^n)^- \in F$.
- (ii) $(x \odot y)^\sim \in F$, implies $(x^n)^\sim \in F$ or $(y^n)^\sim \in F$.

Proof. (i) Let F be an n -fold obstinate filter of A and $(x \odot y)^- \in F$. Since F is a proper filter, $x \odot y \notin F$. Then by (F2), $x \notin F$ or $y \notin F$. By Proposition 3.6(ii), $(x^n)^- \in F$ and $(y^n)^- \in F$.

(ii) The proof is similar to (i).

Lemma 3.1 (i) *Let $\varphi \in \text{Hom}(A, B)$ and G be an n -fold obstinate filter of B . Then the inverse image of G is an n -fold obstinate filter of A .*

(ii) *Let $\varphi : A \rightarrow B$ be a pseudo-hoop isomorphism and $F \in \mathcal{F}(A)$ be an n -fold obstinate filter. Then $\varphi(F)$ is an n -fold obstinate filter of B .*

(iii) *Let $\varphi : A \rightarrow B$ be a pseudo-hoop surjective and A be an n -fold obstinate pseudo-hoop. Then B is an n -fold obstinate pseudo-hoop.*

Proof. (i) Let G be an n -fold obstinate filter of B and $x \in A$ but $x \notin \varphi^{-1}(G)$. Then $\varphi(x) \notin G$, and so by Proposition 3.6(ii), $((\varphi(x))^n)^-, ((\varphi(x))^n)^\sim \in G$. By Definition 2.7, we have $\varphi((x^n)^-), \varphi((x^n)^\sim) \in G$. Then $(x^n)^-, (x^n)^\sim \in \varphi^{-1}(G)$. Therefore, $\varphi^{-1}(G)$ is an n -fold obstinate filter of A .

(ii) It is easy to see that, if $F \in \mathcal{F}(A)$, since φ is a pseudo-hoop isomorphism, then $\varphi(F) \in \mathcal{F}(B)$. Now, let $y_1, y_2 \notin \varphi(F)$. Then $\varphi^{-1}(y_1), \varphi^{-1}(y_2) \notin F$. Since F is an n -fold obstinate filter, then $\varphi^{-1}((y_1)^n \rightarrow y_2) = (\varphi^{-1}(y_1))^n \rightarrow \varphi^{-1}(y_2) \in F$ and so $(y_1)^n \rightarrow y_2 \in \varphi(F)$. By the similar way, we can get that $(y_2)^n \rightarrow y_1, (y_1)^n \rightsquigarrow y_2, (y_2)^n \rightsquigarrow$

$y_1 \in \varphi(F)$. Therefore, $\varphi(F)$ is an n -fold obstinate filter of B .

(iii) Let $y \in B$. Then there exists $x \in A$ such that $y = \varphi(x)$ and so $y^n = \varphi(x^n) = \varphi(0) = 0$. Therefore, B is an n -fold obstinate pseudo-hoop.

Theorem 3.3 *The following conditions are equivalent:*

- (i) any filter $F \in \mathcal{F}(A)$ is an n -fold obstinate filter of A ,
- (ii) $\{1\}$ is an n -fold obstinate filter of A ,
- (iii) A is an n -fold obstinate pseudo-hoop.

Proof. (i) \Rightarrow (ii) The proof is clear.

(ii) \Rightarrow (i) By Proposition 3.9, the proof is clear.

(ii) \Rightarrow (iii) Since $A \cong A/\{1\}$ and $\{1\}$ is an n -fold obstinate filter, then by Theorem 3.2 and Lemma 3.1(iii), A is an n -fold obstinate pseudo-hoop.

(iii) \Rightarrow (ii) Let A be an n -fold obstinate pseudo-hoop and $1 \neq x \in A$. Since, $x^n = 0$, by Proposition 2.2(iii), $(x^n)^- = (x^n)^\sim = 1 \in \{1\}$. Then by Proposition 3.6(ii), $\{1\}$ is an n -fold obstinate filter of A .

Proposition 3.11 *Let F be an n -fold obstinate filter of A . Then the following conditions are hold:*

- (i) $[F \cup G]$ is an n -fold obstinate filter of A , for any $G \in \mathcal{F}(A)$.
- (ii) $F(x)$ is an n -fold obstinate filter of A , for all $x \in A$.

Proof. (i) Let $x \notin [F \cup G]$. Then $x \notin F$ and $x \notin G$. By Proposition 3.6(ii), $(x^n)^- \in F$. Thus $(x^n)^- \in [F \cup G]$. By Proposition 3.6(ii), $[F \cup G]$ is an n -fold obstinate filter of A .

(ii) We consider the following cases:

Case 1: If $x \in F$, then $F(x) = F$.

Case 2: If $x \notin F$ and $y \notin F(x)$, $y \neq x$, then $y \notin F$ and by Proposition 3.6(ii), $(y^n)^- \in F$. Hence $(y^n)^- \in F(x)$. By Proposition 3.6(ii), $F(x)$ is an n -fold obstinate filter of A .

4 Relation between n -fold filters in pseudo-hoops

In this section, we investigate the relationship between n -fold obstinate filters and other filters and

n -fold filters in pseudo-hoops.

Theorem 4.1 Every n -fold obstinate filter of A is a maximal filter of A .

Proof. Let F be an n -fold obstinate filter of A which is not a maximal filter of A . Then there exists a proper filter G of A such that $F \subseteq G$. Let $x \in G \setminus F$. Since F is an n -fold obstinate filter, by Proposition 3.6(ii), $(x^n)^- \in F$. Since $(x^n)^- \in G$ and $x^n \in G$, by Proposition 2.2(v), $x^n \odot (x^n)^- = 0 \in G$, which is a contradiction. Therefore, F is a maximal filter.

The next example shows that the converse of Theorem 4.1, is not true, in general.

Example 4.1 Let $(A = \{0, a, b, c, d, 1\}, \leq)$ be a poset. Define operations \odot, \rightsquigarrow and \rightarrow on A as follows,

$\rightarrow, \rightsquigarrow$	0	a	b	c	d	1
0	1	1	1	1	1	1
a	c	1	b	c	b	1
b	d	a	1	b	a	1
c	a	a	1	1	a	1
d	b	1	1	b	1	1
1	0	a	b	c	d	1

\odot	0	a	b	c	d	1
0	0	0	0	0	0	0
a	0	a	d	0	d	a
b	0	d	c	c	0	b
c	0	0	c	c	0	c
d	0	d	0	0	0	d
1	0	a	b	c	d	1

By routine calculations, we can see that $(A, \odot, \rightarrow, \rightsquigarrow, 0, 1)$ is a bounded pseudo-hoop. It is clear that $F = \{1, a\}$ is a maximal filter but it is not an 1-fold obstinate filter. Because $b \notin F$ and $b^- = d \notin F$.

Corollary 4.1 Every n -fold obstinate filter of pseudo \vee -hoop A is a prime filter of A .

Proof. By Theorem 4.1 and Proposition 2.5, the proof is clear.

Proposition 4.1 Any 1-fold obstinate filter F is a normal filter of A .

Proof. Let F be an 1-fold obstinate filter and $x \rightarrow y \in F$. We consider the following cases:

Case 1: If $y \in F$, then by Proposition 2.1(iii), $y \leq x \rightsquigarrow y$. By (F1), $x \rightsquigarrow y \in F$.

Case 2: If $x, y \notin F$, then by assumption, $x \rightsquigarrow y \in F$.

Case 3: If $x \in F$, then by Proposition 2.1(v), $(x \rightarrow y) \odot x \leq y$. By (F1) and (F2), $y \in F$. Hence by Case 1, $x \rightsquigarrow y \in F$.

Therefore, F is a normal filter of A .

In the following example we show that the converse of Proposition 4.1, is not true, in general.

Example 4.2 In Example 4.1, $F = \{1\}$, is a normal filter but it is not an n -fold obstinate filter. Because, $a^n \rightarrow b = b$ and $b^n \rightarrow a = a \notin F$.

Theorem 4.2 Let F be an n -fold obstinate filter of A . Then F is an n -fold implicative filter.

Proof. Assume that F is not an n -fold implicative filter. Then there exist $x, y \in A$, such that $1 \rightarrow ((x^n \rightarrow y) \rightsquigarrow x) \in F$ but $x \notin F$. By Proposition 2.3(ii), $(x^n \rightarrow y) \rightsquigarrow x \in F$. We consider two cases:

Case 1: If $y \in F$, then since $y \leq x^n \rightarrow y$, so by (F1), $x^n \rightarrow y \in F$. By Proposition 2.3(iii), since $(x^n \rightarrow y) \rightsquigarrow x \in F$ and $x^n \rightarrow y \in F$, we get, $x \in F$, which is a contradiction.

Case 2: If $y \notin F$, then since F is an n -fold obstinate filter, $x^n \rightarrow y \in F$. By Proposition 2.3(iii), since $(x^n \rightarrow y) \rightsquigarrow x \in F$ and $x^n \rightarrow y \in F$, we get, $x \in F$, which is a contradiction.

Therefore, F is an n -fold implicative filter of A .

Lemma 4.1 Any filter F of A is an n -fold positive implicative filter if and only if for all $x \in A$, $F_x = \{y \in A \mid x^n \rightarrow y \text{ and } x^n \rightsquigarrow y \in F\}$ is a filter of A .

Proof. Let F be an n -fold positive implicative filter of A . Since $x^n \rightarrow 1 = 1 \in F$, we have $1 \in F_x$. Let $y, z \in A$ such that $y, y \rightarrow z \in F_x$. Then $x^n \rightsquigarrow y \in F$ and $x^n \rightarrow (y \rightarrow z) \in F$. Thus $x^n \rightarrow z \in F$, and so $z \in F_x$. Therefore, F_x is a filter of A .

Conversely, suppose F_x is a filter of A , for all $x \in A$. Let $x, y, z \in A$ such that $x^n \rightarrow (y \rightarrow z) \in F$ and $x^n \rightsquigarrow y \in F$. Then $y, y \rightarrow z \in F_x$. Thus $z \in F_x$, and so $x^n \rightarrow z \in F$. The proof of other cases is similar, too.

Theorem 4.3 *If F is a maximal and n -fold positive implicative filter of A , then F is an n -fold obstinate filter of A .*

Proof. Let F be a maximal and n -fold positive implicative filter of A and $x, y \in A \setminus F$. Then by Lemma 4.1, $F_x = \{b \in A \mid x^n \rightarrow b \text{ and } x^n \rightsquigarrow b \in F\}$ and $F_y = \{b \in A \mid y^n \rightarrow b, y^n \rightsquigarrow b \in F\}$ are filters of A .

Let $z \in F$. Then by Proposition 2.1(iii), $z \leq x^n \rightarrow z$ and by (F1), $x^n \rightarrow z \in F$. Thus, $z \in F_x$ and so $F \subseteq F_x$. On the other hand, $x^n \rightarrow x = 1 \in F$, so $x \in F_x$. By assumption, $x \notin F$. Hence $F \subsetneq F_x \subseteq A$. Since F is a maximal filter of A , $F_x = A$. Hence $y \in F_x$ or equivalently $x^n \rightarrow y \in F$. Similarly $x^n \rightsquigarrow y \in F$, $y^n \rightarrow x \in F$ and $y^n \rightsquigarrow x \in F$.

Proposition 4.2 [2] *Let F be a normal filter of A .*

- (i) *If for all $x \in A$, $x^n \rightarrow x^{2n} \in F$ or $x^n \rightsquigarrow x^{2n} \in F$, then F is an n -fold positive implicative filter of A .*
- (ii) *If F is an n -fold implicative filter of A , then F is an n -fold fantastic filter of A .*
- (iii) *$\{1\}$ is an n -fold fantastic filter, if and only if A is an n -fold fantastic pseudo-hoop.*

Theorem 4.4 *Let A be a pseudo \vee -hoop. Then F is an n -fold obstinate filter if and only if F is a prime and n -fold implicative filter.*

Proof. If F is an n -fold obstinate filter, then by Corollary 4.1 and Theorem 4.2, the proof is clear.

Conversely, assume that F is a prime filter and n -fold implicative filter of A such that $x \in A \setminus F$. We show that $x \vee (x^n)^- \in F$ and $x \vee (x^n)^\sim \in F$, for all $x \in A$. Since F is an n -fold implicative filter, if $(x^n)^\sim \rightarrow x \in F$ then $x \in F$. Also $(x^n)^- \rightsquigarrow x \in F$ implies $x \in F$. Now, we must show that $t = x \vee (x^n)^- \in F$. Since $x \leq t$, we have $x^n \leq t^n$ and then by Proposition 2.2(i), $(t^n)^- \leq (x^n)^- \leq (x^n)^- \vee x = t$. So $(t^n)^- \rightsquigarrow t = 1 \in F$. Hence, we get that $t \in F$. The other case is similar. Thus $x \vee (x^n)^- \in F$. Since F is a prime filter and $x \notin F$, we have $(x^n)^- \in F$. Therefore, F is an n -fold obstinate filter of A .

Proposition 4.3 *Let F be a normal n -fold obstinate filter of A . Then:*

- (i) *F is an n -fold positive implicative filter,*
- (ii) *F is an n -fold fantastic filter.*

Proof. (i) We consider two cases:

Case 1: Let $x \in F$. Then by (F2), $x^{2n} \in F$ and by Proposition 2.1(iii), $x^{2n} \leq x^n \rightarrow x^{2n}$. By (F1), $x^n \rightarrow x^{2n} \in F$.

Case 2: Let $x \notin F$. Then by assume $(x^n)^- \in F$. By Proposition 2.2(vi), $(x^n)^- \leq x^n \rightarrow x^{2n}$ and by (F1), $x^n \rightarrow x^{2n} \in F$. Therefore, by Proposition 4.2(i), F is an n -fold positive implicative filter of A .

(ii) By Theorems 4.2 and 4.2(ii), F is an n -fold fantastic filter of A .

Theorem 4.5 (i) *If F is an n -fold fantastic filter of A , then $((x^n)^-)^{\sim} \rightarrow x \in F$ and $((x^n)^{\sim})^- \rightarrow x \in F$.*

(ii) *If $D_e(A) = \{x \in A \mid x^- = x^\sim = 0\} = A \setminus \{0\}$, then every n -fold fantastic filter is an n -fold obstinate filter of A .*

(iii) *Let F be an n -fold fantastic filter and for all $x, y \in A$, if $(x^n \odot y^n)^- \in F$, then $(x^n)^- \in F$ or $(y^n)^- \in F$. Also, $(x^n \odot y^n)^\sim \in F$ implies $(x^n)^\sim \in F$ or $(y^n)^\sim \in F$. Then F is an n -fold obstinate filter of A .*

Proof. (i) Since $0 \rightarrow x = 0 \rightsquigarrow x = 1 \in F$ and F is an n -fold fantastic filter, then $((x^n)^-)^{\sim} \rightarrow x \in F$ and $((x^n)^{\sim})^- \rightarrow x \in F$.

(ii) Let F be a proper n -fold fantastic filter of A . Then $0 \notin F$, and so $(x^n)^-$, $(x^n)^\sim \notin F$, for any $0 \neq x \in A$ and $n \geq 1$. By assumption and (i), $((x^n \rightarrow 0) \rightsquigarrow 0) \rightarrow x = (0 \rightsquigarrow 0) \rightarrow x = 1 \rightarrow x = x \in F$. Hence, by Proposition 3.6(ii), F is an n -fold obstinate filter.

(iii) Assume F is an n -fold fantastic filter of A such that $x \notin F$. It is enough to prove that $(x^n)^-, (x^n)^\sim \in F$. Let $(x^n)^- \notin F$, by the contrary. Then by Proposition 2.2(v), $(x^n \odot (x^n)^-)^{\sim} = 0^\sim = 1 \in F$. By assumption $((x^n)^-)^{\sim} \in F$. Since F is an n -fold fantastic filter, by (i), $((x^n)^-)^{\sim} \rightarrow x \in F$. By Proposition 2.3(ii), $x \in F$, which is a contradiction. Hence, $(x^n)^\sim \in F$. By the similar way, we get that $(x^n)^- \in F$. Therefore, F is an n -fold obstinate filter of A .

Proposition 4.4 *Let A be an n -fold fantastic pseudo-hoop and if for all $x, y \in A$, $x^n \odot y^n = 0$*

implies $x^n = 0$ or $y^n = 0$. Then A is an n -fold obstinate pseudo-hoop.

Proof. If A is an n -fold fantastic pseudo-hoop, then by Proposition 4.2(iii), $\{1\}$ is an n -fold fantastic filter of A . By hypothesis and Theorem 4.5(iii), $\{1\}$ is an n -fold obstinate filter of A and so by Theorem 3.3(iii), A is an n -fold obstinate pseudo-hoop.

Proposition 4.5 Let A be a bounded simple pseudo-hoop. Then A is an n -fold obstinate pseudo-hoop, for some $n \in \mathbb{N}$.

Proof. If $1 \neq x \in A$, then $[x] = A$ and so $0 \in [x]$. Hence for some $m \in \mathbb{N}$, $x^m = 0$. Let $n = \max\{m \mid x \in A\}$. Then A is an n -fold obstinate pseudo-hoop.

Theorem 4.6 Let F be an n -fold obstinate filter of A . Then A/F is a local and simple pseudo-hoop.

Proof. Let F be an n -fold obstinate filter of A . Then by Theorem 4.1, F is a maximal filter of A and so A/F is a local and simple pseudo-hoop.

Notation: A partially ordered set (P, \leq) is called to be of the finite length if the length of all chains in P are finite.

Theorem 4.7 Let A be a pseudo-hoop of finite length. Then there exists $n \in \mathbb{N}$ such that every maximal filter of A is an n -fold obstinate filters of A .

Proof. Let n be the length of the greatest chain in A . Then by Theorem 4.1, every n -fold obstinate filter of A is a maximal one. Now, let $F \in \text{Max}(A)$. Then, we show that F is an n -fold obstinate filter. Assume $x \notin F$. Since F is a maximal filter of A , by Proposition 2.4, then $(x^t)^- \in F$, for some $t \in \mathbb{N}$. If $t \leq n$, then by Proposition 2.1(v), $x^n \leq x^t$, so by Proposition 2.2(i), $(x^t)^- \leq (x^n)^-$. By (F1), $(x^n)^- \in F$. Let $n < t$. Since $0 \leq x^n \leq x^{n-1} \leq \dots \leq x^2 \leq x \leq 1$ and A is finite length. Then by assumption, there is a $s \in \{1, 2, \dots, n\}$ such that $x^s = x^{s+1}$, so $x^n = x^t$. It follows that $(x^n)^- \in F$. Therefore, F is an n -fold obstinate of A .

Theorem 4.8 Let A be an n -fold obstinate pseudo-hoop. Then the following conditions are hold:

- (i) A is an n -fold fantastic pseudo-hoop,
- (ii) A is an n -fold positive implicative pseudo-hoop,
- (iii) A is an n -fold implicative pseudo-hoop,
- (iv) A is a local pseudo-hoop,
- (v) A is a simple pseudo-hoop.

Proof. (i) Let A be an n -fold obstinate pseudo-hoop. Then by Theorem 3.3(ii), $\{1\}$ is an n -fold obstinate filter of A . By Proposition 4.3(ii), $\{1\}$ is an n -fold fantastic filter of A . then by Proposition 4.2(iii), A is an n -fold fantastic pseudo-hoop.

(ii) Let A be an n -fold obstinate pseudo-hoop. Then $x^n = 0$, and so $x^{n+1} = x^n$. Hence, A is an n -fold positive implicative pseudo-hoop.

(iii) Let A be an n -fold obstinate pseudo-hoop. Then by Proposition 2.1(ii), $(x^n \rightarrow 0) \rightsquigarrow x = 1 \rightsquigarrow x = x$ and $(x^n \rightsquigarrow 0) \rightarrow x = 1 \rightarrow x = x$. Therefore, A is an n -fold implicative pseudo-hoop.

(iv) Since for any $1 \neq x \in A$, $x^n = 0$, then $\text{ord}(x) < \infty$. Hence, A is a local pseudo-hoop.

(v) Let A be an n -fold obstinate pseudo-hoop and $1 \neq x \in F$. Then by (F2), $0 = x^n \in F$. Therefore, A is a simple pseudo-hoop.

In the following diagram, we show the relationship between n -fold obstinate filter and other filters of pseudo-hoop, where the condition (*) is $x^n \odot y^n = 0 \Rightarrow x^n = 0$ or $y^n = 0$.

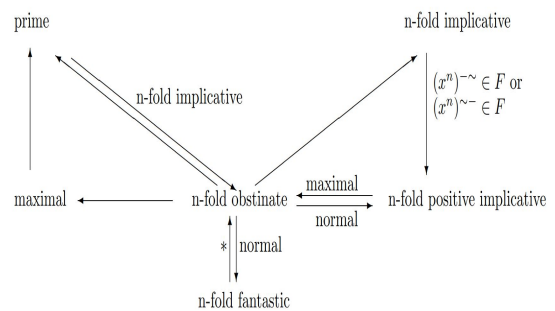


Figure 1: First-type nanostar dendrimer, $NS_1[2]$

5 Conclusion

In this paper, we have considered the folding theory of a filter which is a generalization of a filter in pseudo-hoop. We have provided conditions for a filter to be an n -fold obstinate filter of a pseudo-hoop. So we discuss on concept n -fold obstinate pseudo-hoops. Then we studied relationships between n -fold obstinate pseudo-hoops and some other special pseudo-hoops, such as simple pseudo-hoop and local pseudo-hoop. On the other hands, we introduced the notion of an n -fold obstinate filter in pseudo-hoop. Then we studied relationships between an n -fold obstinate filter and some other special n -fold filter, such as n -fold fantastic, n -fold positive implicative and n -fold implicative filter.

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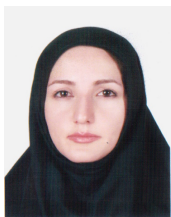
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